

MONOTONIC ITERATIVE ALGORITHMS FOR A QUASICOMPLEMENTARITY PROBLEM*

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Abstract

We present two iterative algorithms, so called *SCP* and *SA* respectively, for solving quasicomplementarity problem (*QCP*). Algorithm *SCP* is to approximate *QCP* by a sequence of ordinary complementarity problems (*CP*). *SA* is a Schwarz algorithm which can be implemented parallelly. We prove the algorithms above are monotonically convergent.

Key words: Quasicomplementarity problem, Iterative algorithm, Monotonic convergence, Schwarz algorithm.

1. Introduction

Consider the following *QCP*: find $u \in R^n$ such that

$$\min\{Au - f, \quad u - Bu\} = 0, \quad (1)$$

where $A, B: R^n \rightarrow R^n$ are operators, $f \in R^n$. The quasivariational inequality which is equivalent to (1) sounds: to find $u \in R^n$ such that $u \geq Bu$ and

$$(Au, v - u) \geq (f, v - u), \quad \forall v \geq Bu. \quad (2)$$

If $Bu \equiv c \in R^n$ for any $v \in R^n$ then *QCP*(1) reduces into *CP*. (1) appears in mathematical programming (see, for example, [2] and the references therein), also comes from the discretization of *QCP* in mathematical physics and control theory (see [1]). For various generalization, see [3] and the references therein.

We assume in this paper that A is a strictly T-monotonic operator, that is:

$$(Au - Av, (u - v)^+) \geq 0, \quad \forall u, v \in R^n,$$

where the equality holds only if $(u - v)^+ = 0$, $v^+ = \max\{v, 0\}$. The examples in [6] show that strictly T-monotonic operator is widely applicable.

We propose two algorithms for solving (1). The first (*SCP*) is to solve iteratively a sequence of *CP*, which produces a sequence of approximate solutions convergent monotonically to a solution of (1). The second (*SA*) is a Schwarz algorithm, which is parallel algorithm and produce super(sub)solution sequence of (1), convergent monotonically to a solution of (1).

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Schwarz algorithms have been developed rapidly (for example, see [4], [6] and the references therein), But Algorithm SA is the first Schwarz algorithm for quasivariational inequality.

2. Sequential CP Algorithm

Algorithm SCP

1⁰. Take $u^0 \in R^n$, $k := 0$;

2⁰. Find $u^{k+1} \in R^n$ such that

$$\min \{Au^{k+1} - f, u^{k+1} - Bu^k\} = 0; \quad (3)$$

3⁰. $k := k + 1$, go to 2⁰.

At first we study an operator F related to (3) and defined as follows: for any $x \in R^n$ define $y = Fx$ as the solution of the following CP:

$$\min\{Ay - f, y - Bx\} = 0. \quad (4)$$

It has a unique solution if operator A is continuous, strictly T-monotonic and coercive in some sense (see, for example, [5]). Then F is well-defined.

We call an operator B order-preserved if $v \leq w$ implies $Bv \leq Bw$.

Lemma 1. Assume B is continuous, order-preserved and there exists $b \in R^n$ such that

$$Bv \leq b, \quad \forall v \in R^n. \quad (5)$$

Assume A is continuous, strictly T-monotonic and

$$\frac{(Av, v - b)}{\|v\|} \rightarrow +\infty \quad (\|v\| \rightarrow \infty). \quad (6)$$

Then the range of F , denoted by $R(F)$, is bounded. That is, there exist $p, q \in R^n$ such that

$$p \leq Fx \leq q, \quad \forall x \in R^n.$$

Proof. Since (6) implies the coerciveness condition of A in [5], (4) has a unique solution for any $x \in R^n$ and F is well-defined. Hence we have

$$\min\{AFx - f, Fx - Bx\} = 0, \quad \forall x \in R^n.$$

Then for any $x \in R^n$ we have

$$(AFx, v - Fx) \geq (f, v - Fx), \quad \forall v \geq Bx.$$

Letting $v = b$ in it we obtain

$$(AFx, Fx - b) \leq (f, Fx - b), \quad \forall x \in R^n$$

which combining with (6) yields that F is bounded.

Remark 1. If there exists a constant α such that

$$(Av - Aw, v - w) \geq \alpha \|v - w\|^2, \quad \forall v, w \in R^n$$